## MUTHAYAMMAL ENGINEERING COLLEGE

## (An Autonomous Institution)

(Approved by AICTE, New Delhi, Accredited by NAAC \& Affiliated to Anna University)
Rasipuram - 637 408, Namakkal Dist., Tamil Nadu
MUST KNOW CONCEPTS
B.E - CY

Course Code \& Course Name : 19CYC01 \& DATA STRUCTURES
Year/Sem/Sec : II/III/-

| Subject |  | 19CYC01 \& DATA STRUCTURES |  |  |
| :---: | :---: | :---: | :---: | :---: |
| S.No | Term | Notation <br> (Symbol) | Concept/Definition/Meaning/Units/Equa tion/Expression | Units |
| Unit-I : Introduction and List |  |  |  |  |
| 1 | Data |  | Data are simply values or sets of values |  |
| 2 | Information |  | Processed Data |  |
| 3 | Datum |  | Singular form of Data |  |
| 4 | Data |  | Plural form of Data |  |
| 5 | Data structures | DS | Way of organizing data in a computer called DS |  |
| 6 | Classification of DS | ST | Static data structures <br> Dynamic data structures |  |
| 7 | Static data structures | $-51$ | Fixed size data structure. EX: Array, pointers, structures |  |
| 8 | Dynamic data structures |  | Variable size data structure. Ex: linked lists, stacks, queues, trees |  |
| 9 | Types of data structure |  | Linear data structure. <br> Non-linear data structure |  |
| 10 | Linear data structures |  | Data are arranged in sequential order |  |
| 11 | Non-linear data structure |  | Data structures that don't have a linear relationship between its adjacent elements but have a hierarchical relationship |  |
| 12 | Abstract Data Type | ADT | Set of operations for which the implementation of the data structure is not specified |  |



| 31 | Pop operation |  | Removing an element from stack |  |
| :---: | :---: | :---: | :---: | :---: |
| 32 | Peek operation |  | Viewing top element of stack |  |
| 33 | Empty stack |  | If top=-1 represent empty stack |  |
| 34 | Full |  | If top=maxsize-1 represent full stack |  |
| 35 | Queue |  | Queue is an ordered collection of elements in which insertions and deletions take place in 2 ends |  |
| 36 | Rear end |  | The end from which elements are added referred to rear end |  |
| 37 | Front end |  | End from which deletions are made is referred to as the front end |  |
| 38 | Priority queue |  | Priority queue is a collection of elements, each containing a key referred as the priority for that element |  |
| 39 | Enqueue |  | Inserting an element in queue |  |
| 40 | Dequeue |  | Removing an element from queue |  |
| 41 | Front |  | Ptr points to 1,st element of queue |  |
| 42 | Rear |  | Ptr points to last element of queue |  |
| 43 | Types of queues |  | - Linear queues <br> - Circular queues <br> - Priority queue |  |
| 44 | Applications of stacks |  | - Reversing a string <br> - Balanced parenthesis <br> - Evaluation of arithmetic expressions |  |
| 45 | Underflow |  | Checking queue is empty (contain no elements in array) called underflow |  |
| 46 | Overflow |  | Checking queue is full (contain all elements in array) called overflow |  |
| 47 | LIFO |  | Last in first out (principle followed by stack) |  |
| 48 | FIFO |  | First in first out ( principle followed by stack queue) |  |
| 49 | Max heap |  | The key at root must be maximum among all keys present in binary heap |  |
| 50 | Min heap |  | The key at root must be minimum among all keys present in binary heap |  |


| Unit-III : Tree and Binary Search Tree |  |  |  |  |
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| 51 | Tree |  | $\begin{array}{l}\text { A tree is a non-linear data structure, which } \\ \text { represents hierarchical relationship } \\ \text { between individual data items }\end{array}$ |  |
| 52 | Height of a Tree |  | $\begin{array}{l}\text { Length of the longest path from the root to } \\ \text { a leaf }\end{array}$ |  |
| 53 | Path in a tree | $\begin{array}{l}\text { Sequence of distinct nodes in which } \\ \text { successive nodes are connected by edges }\end{array}$ |  |  |
| 54 | Leaf node | A node that has no children |  |  |\(\left.] \begin{array}{l}A binary tree is a tree in which every non- <br>

leaf node has at most two children\end{array}\right]\)

| 72 | Fields of Double linked list node |  | Data, next and previous |  |
| :---: | :---: | :---: | :---: | :---: |
| 73 | previous |  | Address of previous node of list |  |
| 74 | Isempty of list () |  | If head== NULL represent empty list |  |
| 75 | Traversing |  | Operation perform viewing of all element in the list |  |
| Unit-IV : Graphs |  |  |  |  |
| 76 | Graph |  | A graph is a non-linear data structure that represents less relationship between its adjacent elements. There is no hierarchical relationship between the adjacent elements in case of graphs |  |
| 77 | Undirected graph |  | If an edge between any two nodes in a graph is not directionally oriented a graph is called as undirected graph |  |
| 78 | Directed graph |  | If an edge between any two nodes in a graph is directionally oriented, a graph is called as directed graph; it is also referred as a digraph |  |
| 79 | Cycle |  | A cycle is a path containing at least three vertices such that the starting and the ending vertices are the same. |  |
| 80 | Weighted graph |  | A graph is said to be weighted graph if every edge in the graph is assigned some weight or value |  |
| 81 | Minimum spanning trees |  | A minimum spanning tree is one of the spanning trees of the graph which has the smallest sum of weights amongst all spanning trees. |  |
| 82 | DFS |  | DFS means Depth First search it is like a preorder traversal of a tree. It is continuous searching for the unvisited nodes in the forward direction based on the recursive process |  |
| 83 | Complete Graph |  | In a graph if there exists the path from any vertex to any other vertex, then the graph is called as Complete Graph |  |
| 84 | BFS |  | BFS performs simultaneous exploration starting from a common point and spreading out independently |  |
| 85 | Self loop |  | In graph theory, a loop is an edge that connects a vertex to itself |  |
| 86 | Representation of Graph |  | - Adjacency List <br> - Adjacency Matrix |  |
| 87 | Data Structure used in BFS |  | Queue |  |


| 88 | Data Structure used in DFS |  | Stack |  |
| :---: | :---: | :---: | :---: | :---: |
| 89 | Vertex |  | Each node of the graph is termed as vertex |  |
| 90 | Edge |  | Edge represents a path between two vertices |  |
| 91 | Adjacency |  | Two nodes or vertices are adjacent if they are connected to each other |  |
| 92 | Path |  | Path represents the series of edges between two vertices |  |
| 93 | Basic operations on the graph |  | - Add vertex <br> - Add Edge <br> - Display Vertex |  |
| 94 | Out Degree |  | Number of outgoing vertex |  |
| 95 | In Degree |  | Number of incoming vertex |  |
| 96 | Degree of a graph |  | Number of incident edges |  |
| 97 | Cycle |  | Cycle is a path which starts and ends with a same vertex |  |
| 98 | Connected graph |  | Has all pairs of vertices connected by at least one path |  |
| 99 | Directed Path |  | It is a path of only directed edges |  |
| 100 | Directed Cycle |  | It is a cycle of only directed edges |  |
| Unit-V : Hashing, SearchingAndSorting |  |  |  |  |
| 101 | Hashing |  | Searching technique in $\mathrm{O}(1)$ time complexity |  |
| 102 | Hash function | ESTU | Hash(key Value)=(key value) mod (Table size) |  |
| 103 | Collision in hashing |  | When an element is inserted, it hashes to the same value as an already inserted element, and then it produces collision. |  |
| 104 | Separate chaining |  | Separate chaining is a collision resolution technique to keep the list of all elements that hash to the same value |  |
| 105 | Open addressing |  | Open addressing is a collision resolving strategy in which, if collision occurs alternative cells are tried until an empty cell is found |  |
| 106 | Types of collision resolution strategies in open addressing |  | - Linear probing <br> - Quadratic probing <br> - Double hashing |  |
| 107 | Probing |  | Process of getting next available hash table array cell |  |


| 108 | Linear probing |  | $\begin{aligned} & \text { F(i)=i. } \operatorname{Hi}(\mathrm{x})=(\text { hash }(\mathrm{x})+\mathrm{f}(\mathrm{i})) \text { modtablesize . } \\ & \mathrm{I}=1,2,3,4 \ldots \ldots \end{aligned}$ |
| :---: | :---: | :---: | :---: |
| 109 | Quadratic probing |  | $\begin{aligned} & \mathrm{F}(\mathrm{i})=\mathrm{i} 2 . \operatorname{Hi}(\mathrm{x})=(\mathrm{hash}(\mathrm{x})+\mathrm{f}(\mathrm{i})) \text { modtablesize } \\ & \mathrm{I}=1,2,3,4 \ldots \ldots \end{aligned}$ |
| 110 | Sorting |  | A sorting algorithm is used to rearrange a given array or list elements in ascending or descending order. |
| 111 | Types of internal sorting |  | - Bubble Sort <br> - Insertion Sort <br> - Selection Sort <br> - Quick Sort <br> - Merge Sort <br> - Heap Sort |
| 112 | Classification of sorting |  | Internal sorting and external sorting |
| 113 | Internal sorting |  | internal sorting the data that has to be sorted will be in the main memory |
| 114 | External sorting |  | External sorting it will on disks, outside main memory |
| 115 | Types of external sorting |  | Two-way merge sort , radix sort |
| 116 | Time complexity of bubble sort |  | $\Theta(\mathrm{n})$ |
| 117 | Divide-and-Conquer |  | Divide: Break the given problem into sub problems of same type. <br> Conquer: Recursively solve these sub problems <br> Combine: Appropriately combine the answers |
| 118 | Not a stable sorting algorithm |  | Bubble sort |
| 119 | Not a stable sorting algorithm | ES | Merge sort FUTIRE |
| 120 | $\mathrm{O}(\mathrm{nlogn})$ | $E S t$ | Running merge sort on an array of size $n$ which is already sorted is |
| 121 | $\mathrm{O}(\mathrm{n} \log \mathrm{n})$ ) |  | The time complexity of a quick sort algorithm |
| 122 | Time complexity of insertion sort |  | $\Theta(\mathrm{n})$ |
| 123 | Mod function \% |  | Returns remainder value |
| 124 | 7\%8 |  | 7 |
| 125 | 10\%8 |  | 2 |
| PLACEMENT QUESTIONS |  |  |  |
| 126 | Three times the first of three consecutive |  | Let the three integers be $x, x+2$ and $x+4$. Then, $3 x=2(x+4)+3 \Leftrightarrow x=11$. |



| 136 | Today is Monday. <br> After 61 days, it will be: |  | Each day of the week is repeated after 7 days. <br> So, after 63 days, it will be Monday. <br> $\therefore$ After 61 days, it will be Saturday. |  |
| :---: | :---: | :---: | :---: | :---: |
| 137 | If $6^{\text {th }}$ March, 2005 is Monday,The day of the week on $6^{\text {th }}$ March, 2004 is |  | The year 2004 is a leap year. So, it has 2 odd days. <br> But, Feb 2004 not included because we are calculating from March 2004 to March 2005. So it has 1 odd day only. <br> $\therefore$ The day on $6^{\text {th }}$ March, 2005 will be 1 day beyond the day on $6^{\text {th }}$ March, 2004. Given that, $6^{\text {th }}$ March, 2005 is Monday. <br> $\therefore 6^{\text {th }}$ March, 2004 is Sunday ( 1 day before to $6^{\text {th }}$ March, 2005). |  |
| 138 | The days in $x$ weeks $x$ days? |  | $x$ weeks $x$ days $=(7 x+x)$ days $=8 x$ days. |  |
| 139 | On $8^{\text {th }}$ Feb, 2005 it was Tuesday. The day of the week on $8^{\text {th }} \mathrm{Feb}, 2004$ is |  | The year 2004 is a leap year. It has 2 odd days. <br> $\therefore$ The day on $8^{\text {th }} \mathrm{Feb}$, 2004 is 2 days before the day on $8^{\text {th }} \mathrm{Feb}, 2005$. <br> Hence, this day is Sunday. |  |
| 140 | The greatest number that will divide 43, 91 and 183 so as to leave the same remainder in each case. | $51$ | $\begin{aligned} & \text { Required number }=\text { H.C.F. of }(91-43), \\ & (183-91) \text { and }(183-43) \\ & \quad=\text { H.C.F. of } 48,92 \text { and } 140=4 . \end{aligned}$ |  |
| 141 | The H.C.F. of two numbers is 23 and the other two factors of their L.C.M. are 13 and 14. The larger of the two numbers is: |  | Clearly, the numbers are ( $23 \times 13$ ) and (23 x 14 ). <br> $\therefore$ Larger number $=(23 \times 14)=322$ |  |
| 142 | $\left(112 \times 5^{4}\right)=$ ? |  | $\begin{aligned} & \left(112 \times 5^{4}\right)=112 \times(10) 4=112 \times \\ & 10^{4}=1120000=7000022^{4} 16 \end{aligned}$ |  |
| 143 | It was Sunday on Jan 1,2006 .The day of the week Jan 1, 2010 is |  | On $31^{\text {st }}$ December, 2005 it was Saturday. <br> Number of odd days from the year 2006 to the year $2009=(1+1+2+1)=5$ days. |  |



| 150 | Two trains running in opposite directions cross a man standing on the platform in 27 seconds and 17 seconds respectively and they cross each other in 23 seconds. The ratio of their speeds is: | Let the speeds of the two trains be $x \mathrm{~m} / \mathrm{sec}$ and $\mathrm{y} \mathrm{m} / \mathrm{sec}$ respectively. <br> Then, length of the first train $=27 x$ meters, and length of the second train $=$ $17 y$ meters. $\begin{aligned} & \therefore \frac{27 x+}{x+y}= \\ & \Rightarrow 23 x+17 y=23 x+23 y \\ & \Rightarrow 4 x=6 y \\ & \Rightarrow \\ & \Rightarrow \frac{x}{y}=\frac{3}{2} . \end{aligned}$ |
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## Faculty Team Prepared

## 1. Ms.G.Nivedhitha

## Signatures

